Subfields of \mathbb{R} : NIP, VC Dimension and PAC Learning

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Model Theory of Subfields of $\ensuremath{\mathbb{R}}$



Let \mathcal{L} be any language and let T be a complete \mathcal{L} -theory. A partitioned \mathcal{L} -formula $\varphi(\mathbf{x}; \mathbf{y})$ has the independence property (IP) in T, if for any $n \in \mathbb{N}$, writing $[n] = \{1, \ldots, n\}$, we have



An \mathcal{L} -structure \mathcal{M} has the independence property (IP) if there exists a partitioned \mathcal{L} -formula $\varphi(\mathbf{x}; \mathbf{y})$ that has IP in the theory $\operatorname{Th}_{\mathcal{L}}(\mathcal{M})$.

S. SHELAH, 'Stability, the f.c.p., and superstability; model theoretic properties of formulas in first order theory', *Ann. Math. Logic* **3** (1971) 271–362.

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NIP - not IP

A subfield *K* of \mathbb{R} is called real closed if it is relatively algebraically closed in \mathbb{R} .

Definition.

$$\begin{split} \mathcal{L}_{\mathrm{r}} &:= \{+,-,\cdot,0,1\} - \text{language of rings} \\ \mathcal{L}_{\mathrm{or}} &:= \{+,-,\cdot,0,1,<\} - \text{language of ordered rings} \end{split}$$

Theorem.

A subfield of ${\mathbb R}$ has NIP if it is real closed.

S. SHELAH, 'Strongly dependent theories', Isr. J. Math. 204 (2014) 1–83.

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Conjecture.

A subfield of \mathbb{R} has NIP only if it is real closed.

S. SHELAH, 'Strongly dependent theories', Isr. J. Math. 204 (2014) 1–83.

Subfields of \mathbb{R} with (N)IP



- B. POONEN, 'Uniform first-order definitions in finitely generated fields'. Duke Math. I. **138** (2007) 1–22.
- J. ROBINSON, 'Definability and decision problems in arithmetic', J. Symb. Log. 14 (1949) 98–114.
- J. ROBINSON, 'The undecidability of algebraic rings and fields', Proc. Amer. Math. Soc. 10 (1959) 950-957.

Relation of NIP to Statistical Learning Theory



Given an \mathcal{L} -structure \mathcal{M} and a partitioned \mathcal{L} -formula $\varphi(\mathbf{x}; \mathbf{y})$ with $|\mathbf{x}| = n$ and $|\mathbf{y}| = \ell$, we set

$$\varphi(\mathcal{M}; \boldsymbol{b}) := \{ \boldsymbol{a} \in \mathcal{M}^n \mid \mathcal{M} \models \varphi(\boldsymbol{a}; \boldsymbol{b}) \}$$

for any $\boldsymbol{b} \in M^{\ell}$, and

$$\mathcal{H}_{\varphi} := \{\mathbb{1}_{\varphi(\mathcal{M};b)} \mid b \in M^{\ell}\}.$$

Proposition (Laskowski 1992).

The class \mathcal{H}_{φ} has finite VC dimension if and only if the \mathcal{L} -formula $\varphi(\mathbf{x}; \mathbf{y})$ has NIP in \mathcal{M} .

M. C. LASKOWSKI, 'Vapnik–Chervonenkis classes of definable sets', J. Lond. Math. Soc. 45 (1992) 377–384.

Ingredients of a Learning Problem.

- $\cdot \ \mathcal{X} \mathsf{input} \mathsf{ space}$
- + $\{0,1\}$ output space
- + $\mathcal{Z} = \mathcal{X} \times \{0,1\}$ sample space
- + $\mathcal{H} \subseteq \{0,1\}^{\mathcal{X}} \text{hypothesis space}$

Using an arbitrary distribution \mathbb{D} on $\mathcal{Z} = \mathcal{X} \times \{0, 1\}$, we choose a sequence of iid samples from \mathcal{Z} :

$$\mathbf{z} = ((\mathbf{x}_1, y_1), \ldots, (\mathbf{x}_m, y_m)).$$

These samples provide the input data for a learning algorithm A that determines a hypothesis h = A(z) in H.

S. BEN-DAVID and S. SHALEV-SHWARTZ, Understanding Machine Learning: From Theory to Algorithms, (Cambridge University Press, Cambridge, 2014).

The goal is to minimize the error of h given by

$$\operatorname{er}_{\mathbb{D}}(h) := \mathbb{D}(\{(\mathbf{x}, y) \in \mathcal{Z} \mid h(\mathbf{x}) \neq y\}).$$

More precisely, we want to achieve an error that is close to the approximation error of \mathcal{H} given by

 $\mathsf{opt}_{\mathbb{D}}(\mathcal{H}) := \inf_{h \in \mathcal{H}} \mathsf{er}_{\mathbb{D}}(h).$

S. BEN-DAVID and S. SHALEV-SHWARTZ, Understanding Machine Learning: From Theory to Algorithms, (Cambridge University Press, Cambridge, 2014).

A learning algorithm

$$\mathcal{A}\colon \bigcup_{m\in\mathbb{N}}\mathcal{Z}^m\to\mathcal{H}$$

for \mathcal{H} is said to be probably approximately correct (PAC) if it satisfies the following condition:

$$\begin{aligned} \forall \varepsilon, \delta \in (0,1) \ \exists m_0 = m_0(\varepsilon, \delta) \ \forall m \geq m_0 \ \forall \mathbb{D} : \\ \mathbb{D}^m(\{z \in \mathcal{Z}^m \mid \mathsf{er}_{\mathbb{D}}(\mathcal{A}(z)) - \mathsf{opt}_{\mathbb{D}}(\mathcal{H}) \leq \varepsilon\}) > 1 - \delta. \end{aligned}$$

L. G. VALIANT, 'A Theory of the Learnable', Comm. ACM 27 (1984) 1134–1142.

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The hypothesis class \mathcal{H} is said to be probably approximately correct (PAC) learnable if there exists a learning algorithm for \mathcal{H} that is PAC.

The following result is due to Blumer, Ehrenfeucht, Haussler and Warmuth 1989.

Theorem.

Under certain measurability conditions, a hypothesis class $\mathcal H$ is PAC learnable if and only if its VC dimension is finite.

A. BLUMER, A. EHRENFEUCHT, D. HAUSSLER and M. K. WARMUTH, 'Learnability and the Vapnik-Chervonenkis dimension', J. Assoc. Comput. Mach. **36** (1989) 929–965.

The following result is due to Blumer, Ehrenfeucht, Haussler and Warmuth 1989.

Theorem.

Under certain measurability conditions, a hypothesis class $\mathcal H$ is PAC learnable if and only if its VC dimension is finite.

Corollary.

Given a subfield $K \subseteq \mathbb{R}$ with NIP, any *definable* hypothesis class $\mathcal{H} \subseteq \{0,1\}^{\mathcal{X}}$ with $\mathcal{X} \subseteq K^n$ has finite VC dimension and is thus PAC learnable, provided certain measurability conditions are guaranteed.

Creating a Learning Framework



Assembling the Ingredients.

- $K \subseteq \mathbb{R}$ subfield of \mathbb{R}
- + $\mathcal{L} \in \{\mathcal{L}_{\mathrm{r}}, \mathcal{L}_{\mathrm{or}}\}$ language for model-theoretic examination
- $\mathcal{X} \subseteq K^n definable$ subset of K^n
- $\mathcal{H} \subseteq \{0,1\}^{\mathcal{X}}$ set of hypotheses $h = \mathbb{1}_A : \mathcal{X} \to \{0,1\}$ with *definable* support $A \subseteq \mathcal{X}$

We denote by $\mathcal{B}_{\mathcal{X}}$ and $\mathcal{B}_{\mathcal{Z}}$ the Borel σ -algebras on \mathcal{X} and \mathcal{Z} , respectively.

Lemma.

 $\mathcal{B}_{\mathcal{Z}}=\mathcal{B}_{\mathcal{X}}\otimes\mathcal{P}(\{0,1\}).$

The distributions $\mathbb D$ on $\mathcal Z$ that we consider are all defined on $\mathcal B_{\mathcal Z}.$

Given a hypothesis $h = \mathbb{1}_A$ with definable support $A \subseteq \mathcal{X}$ and a distribution \mathbb{D} defined on $\mathcal{B}_{\mathcal{Z}}$, the error

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\operatorname{er}_{\mathbb{D}}(h) = \mathbb{D}(\{(\mathbf{x}, y) \in \mathcal{Z} \mid h(\mathbf{x}) \neq y\})
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is only well-defined if the support of h is Borel, i.e. $A \in \mathcal{B}_{\mathcal{X}}$.

Question.

Given a subfield $K \subseteq \mathbb{R}$, is any definable set $A \subseteq K^n$ Borel measurable?

Partial Answers.

Yes, if

- A is quantifier-free definable (e.g. if K is real closed),
- A is countable (e.g. if K is countable).

Summary and Future Work

Corollary.

If $K \subseteq \mathbb{R}$ has NIP, then any definable hypothesis class \mathcal{H} has finite VC dimension and is thus PAC learnable, provided certain measurability conditions are guaranteed.

Task.

Identify and study the measurability requirements involved in the Fundamental Theorem.

For instance: Does definability guarantee Borel measurability?

Corollary.

If $K \subseteq \mathbb{R}$ has NIP, then any definable hypothesis class \mathcal{H} is PAC learnable.

Question.

Given $K \subseteq \mathbb{R}$ with IP, does there exist a definable hypothesis class that is not PAC learnable? **Yes!**

Refined Question.

Given $K \subseteq \mathbb{R}$ with IP, does there exist a definable hypothesis class \mathcal{H} generated by a neural network that is not PAC learnable?

Appendix



Measurability – Borel σ -Algebras





Given a hypothesis $h = \mathbb{1}_A$ with support $A \subseteq \mathcal{X}$ and a distribution \mathbb{D} defined on $\mathcal{B}_{\mathcal{Z}}$, the error

$$\operatorname{er}_{\mathbb{D}}(h) = \mathbb{D}(\{(\mathbf{X}, y) \in \mathcal{Z} \mid h(\mathbf{X}) \neq y\}) = \mathbb{D}(\mathcal{Z} \setminus \Gamma(h))$$

is only well-defined if one of the following equivalent conditions is satisfied:

- the map *h* is Borel measurable,
- the support A is Borel, i.e. $A \in \mathcal{B}_{\mathcal{X}}$,
- the graph $\Gamma(h)$ is Borel, i.e. $\Gamma(h) \in \mathcal{B}_{\mathcal{Z}}$.

Problem.

Find sufficient/necessary conditions for the following measurability requirement:

$$\{\mathsf{z}\in\mathcal{Z}^m\mid\mathsf{er}_\mathbb{D}(\mathcal{A}(\mathsf{z}))-\mathsf{opt}_\mathbb{D}(\mathcal{H})\leqarepsilon\}\stackrel{!}{\in}\mathcal{B}^m_\mathcal{Z}$$